

Emerging Application and Algorithm Requirements for Future HPC Systems

July 2008









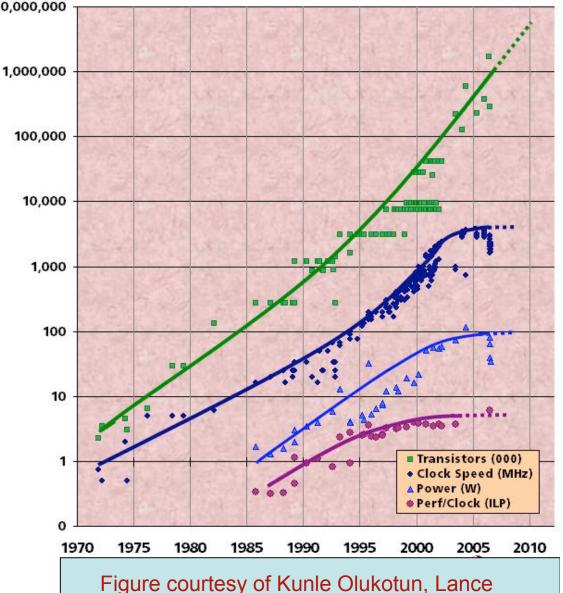
Traditional Sources of Performance Improvement are Flat-Lining

New Constraints

15 years of exponential clock rate growth has ended

But Moore's Law continues!

- How do we use all of those transistors to keep performance increasing at historical rates?
- Industry Response: #cores per chip doubles every 18 months *instead* of clock frequency!



Hammond, Herb Sutter, and Burton Smith

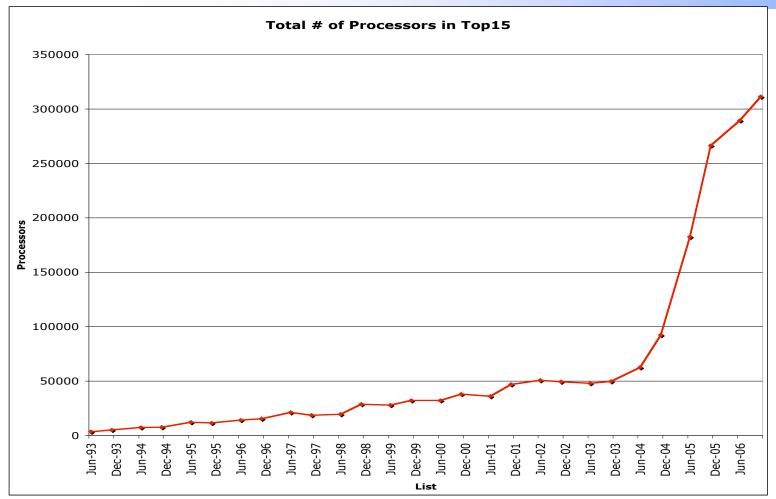




Science

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Growth in HPC System Concurrency



Must ride exponential wave of increasing concurrency for forseeable future!







Application Community's Response to Technology Trends

- Parallel computing has thrived on weak-scaling for past 15 years
- Flat CPU performance increases emphasis on strong-scaling
- Workload Requirements will change accordingly
 - Concurency will increase proportional to system scale (3-5x increase over NERSC-5)
 - Timestepping algorithms will be increasingly driven towards implict or semi-implicit stepping schemes
 - Multiphysics/multiscale problems increasingly rely on spatially adaptive approaches such as Berger-Oliger AMR
 - Strong scaling will push applications towards smaller messages sizes – requiring lighter-weight messaging







NERSC Response To Trends

- Parallel computing has thrived on weak-scaling for past 15 years
- Flat CPU performance increases emphasis on strong-scaling
- NERSC-6 Benchmarks changed accordingly
 - Increased concurrency 4x over NERSC-5 benchmarks
 - Input decks emphasize strong-scaled problems
 - Emphasis on implicit methods
 - New AMR benchmark
 - New UPC benchmark







Materials Science Planewave Density Functional Theory (DFT)

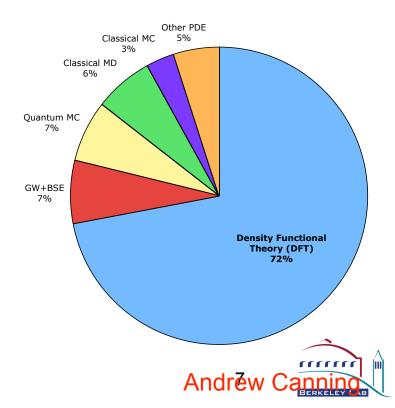






Density Functional Theory (DFT) Algorithm

- Kohn-Sham formalism for computing electronic structure from first principles (DFT Method)
 - Most common implementation is based on expanding the quantum wavefunction into plane-wave (fourier) components
 - This is the method employed by VASP, PARATEC, and Qbox
- Dominant phases of planewave DFT algorithm
 - 3D FFT
 - transforming between real space and reciprocal space
 - O(Natoms²) complexity
 - Subspace Diagonalization
 - O(Natoms³) complexity
 - Orthogonalization
 - dominated by BLAS3
 - ~O(Natoms³) complexity
 - Compute Non-local pseudopotential
 - O(Natoms³) complexity







Future of Materials Science Codes

- For smaller atomic systems (~600-1000 atoms)
 - BLAS dominates at lower concurrencies
 - 3D FFT tends to dominate the computation at high concurrency
 - Due to low computational intensity and small message size (NSF Track-2 bench)
 - Message size can be increased by expending more memory/processor
- For larger atomic systems (>1k atoms), the O(N³) complexity of orthogonalization and computing non-local pseudopotential will dominate
- For O(N³) complexity, moving from teraflops to petaflops only gets you from 1k atoms to 4k atoms.
 - not very impressive given the amount of hardware!
 - Good news is that FLOP rates will be very impressive given increased domination of highly localized BLAS3 operations (eg QBox example)
- For this reason, conventional O(N^3) DFT will be increasingly supplanted by O(N) methods for Petaflop scale calculations!







Anatomy of an O(N) DFT method

(LS3DF as an example)

- Total energy of a system can be decomposed into two parts
 - Quantum mechanical part:
 - wavefunction kinetic energy and exchange correlation energy
 - Highly localized
 - Computationally expensive part to compute
 - Classical electrostatic part:
 - Coulomb energy
 - Involves long-range interactions
 - Solved efficiently using poisson equation even for million atom systems
- LS3DF exploits localization of quantum mechanical part of calculation
 - Divide computational domain into discrete tiles and solve quantum mechanical part
 - Solve global electrostatic part (no decomposition)
 - Very little interprocessor communication required! (almost embarrassingly parallel)
 - Result is O(Natoms) complexity algorithm: enables exploration of larger atomic systems as we move to petaflop and beyond.







Climate



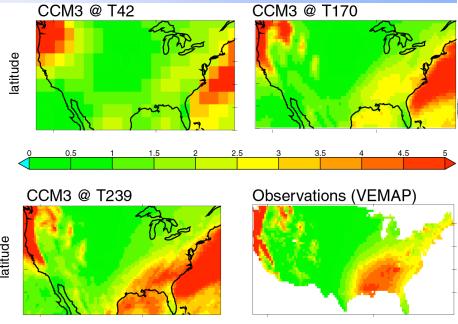




Cloud System Resolving Climate Simulation

- Requires transformational change in science not feasible using current approach
 - The biggest source of climate model errors is poor cloud simulation, especially tropical convection
 - At ~1 km horizontal resolution, cloud systems can be resolved
- DOE Investment in Exascale Computing
 - Climate change is leading justification for general purpose exascale system
 - Not achievable via extrapolation of current approach
 - UN WMO Climate Modeling Summit:
 1km models are the top priority
- Requires substantial code redevelopment to develop cloud-resolving climate model









Global Cloud System Resolving Climate Modeling







Cloud-scale processses
Well understood



Meso-scale statistics Poorly understood



Global scale

This is where parameterization comes in.

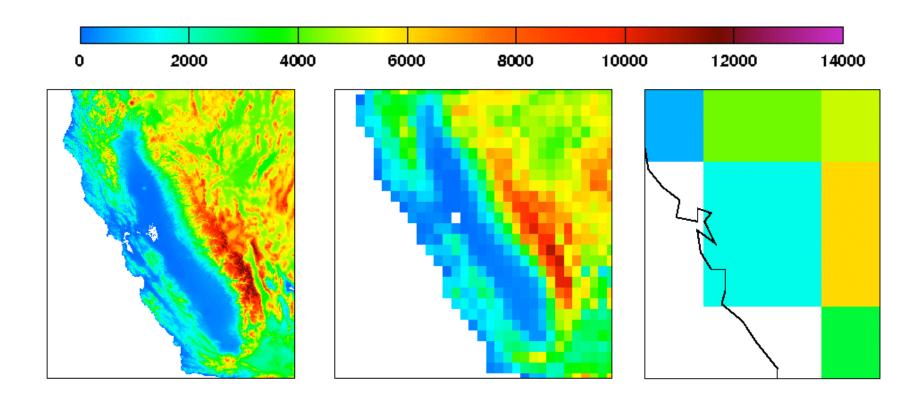
Courtesy Prof. David Randall, Colorado State University

The UN WMO cites the need for Cloud Resolving Models as a Top Priority (cannot be accomplished without 10⁷ improvement in computational capability)





Global Cloud System Resolving Models are a Transformational Change



1km Cloud system resolving models



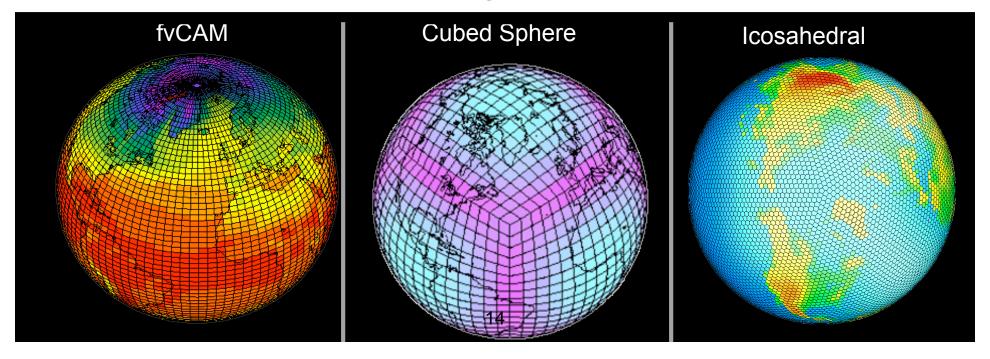
25km
Upper limit of climate models with cloud parameterizations

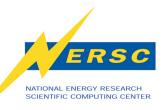
200km
Typical resolution of IPCC AR4 models



Climate Model New Approaches for Massive Parallelism

- Existing Latitude-longitude based algorithm advection algorithm breaks down significantly before 1km scale!
 - Grid cell aspect ratio at the pole is 10000!
 - Advection time step is problematic at this scale
- Ultimately requires new discretization for atmosphere model
 - Must expose sufficient parallelism to exploit power-efficient design
 - Partner with CSU/Randall Group to use the Icosahedral Code
 - Uniform cell aspect ratio across globe

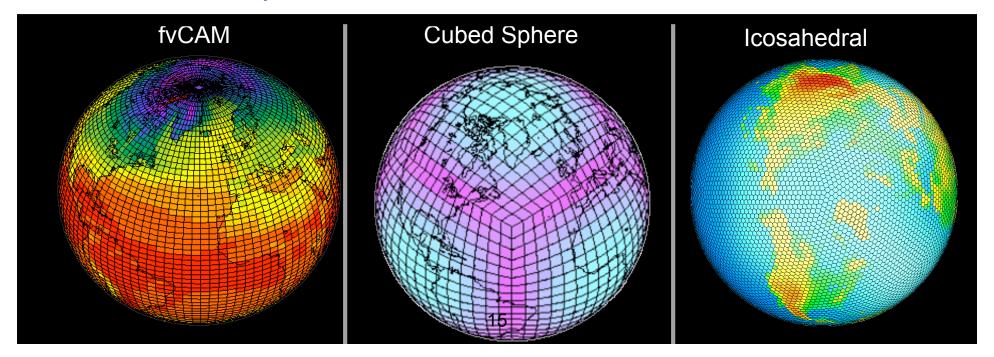




Requirements: 1km Climate Model

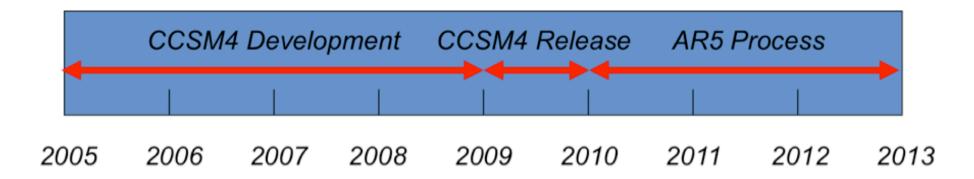
Must maintain 1000x faster than real time for practical climate simulation

- ~2 million horizontal subdomains
- 100 Terabytes of Memory
 - 5MB memory per subdomain
- ~20 million total subdomains
 - 500Mflops sustained per domain
 - Nearest-neighbor communication 250GB/s
- NERSC supports projects developing these new discretizations
 - GFDL Cubed Sphere, CSU Icosahedral model





IPCC AR5 Timeline Coincident with NERSC-6



"The carbon cycle version of CCMS4 will include the additional biogeochemistry, indirect aerosol and land ice components, and the short-term climate simulations will have considerably enhanced atmosphere resolution and, potentially, include the chemistry component. [The] carbon cycle CCSM 4 will be a factor of about five times the CCSM 3 in computing cost. . . . Doing all the proposed IPCC AR 5 runs will stretch the CCSM computing resources to the absolute limit."

Peter R. Gent:

CCSM4 Implementation Plan





Fusion

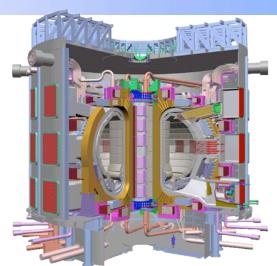






Fusion: Impact of ITER

- Fusion science has been dominated by scaling up first-principles models of specific phenomena
- ITER development requires full-device modeling capability by 2012
 - For shot planning and device control
 - Requires Code-coupling, Multi-scale multiphysics
 - Uncontrolled discharge could damage \$12B device!

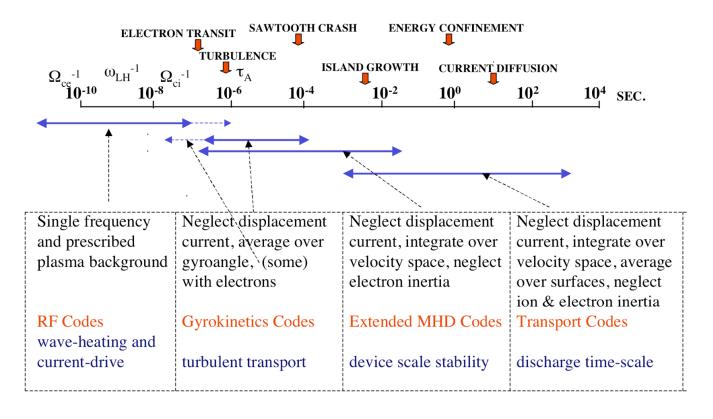


- Requires new code and algorithms to span 12 orders magnitude (Keyes/Jardin)
 - AMR to cover 3 orders of magnitude (time and resolution)
 - Implicit solvers to cover 4 orders magnitude (time)
 - Increased parallel scaling to cover another 3 orders magnitude
 - 2 orders magnitude from higher order elements
- These codes are still in development (and need a platform to support development)
 - SciDAC developing pairwise code coupling

Office of SP will focus on broader coupling for full device modeling capability



Fusion Time and Length Scales

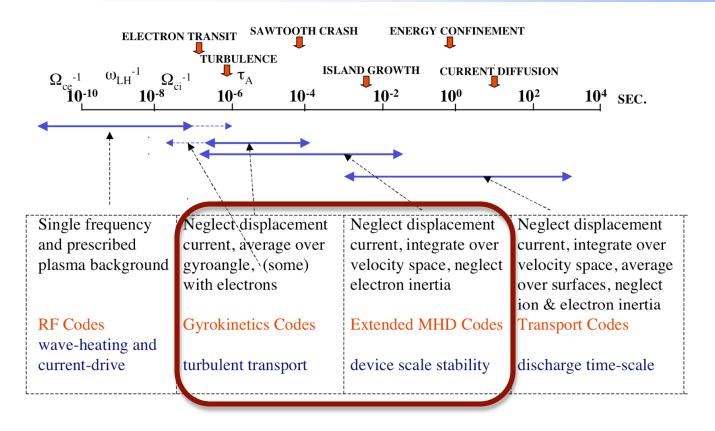








Fusion Time and Length Scales



- Gyrokinetic and MHD codes dominate workload
 - GTC (10%) & GEM (11%) PIC codes dominate Gyrokinetic Codes
 - M3D (10%) & NIMROD (12%) dominate Extended MHD Codes





Emerging Workload Requirements

- Applying computation only where needed
 - AMR: multiscale/multiresolution physics
 - Load balancing issues
 - Locality constraints for prolongation and restriction
 - Many very small (oddly-sized) messages for interconnect
 - Sparse Matrix: Don't compute on non-zeros
 - Very small messages sizes and load balance issues
- Emerging issues with existing applications
 - Implicit Methods
 - Vector inner product required by Krylov subspace algorithms is hampered by latency-bound fast global reductions at massive parallelism
 - Climate Models
 - When science that depends on parameter studies and ensemble runs, capacity and capability are intimately linked!
- I/O Intensive workloads
 - Growth in experimental and sensor data processing

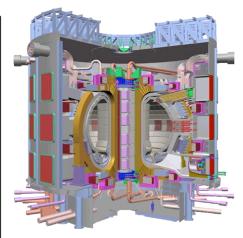






Scaling Fusion Simulations Up to ITER

name	symbol	units	CDX-U	DIII-D	ITER
Field	B_0	Tesla	0.22	1	5.3
Minor radius	а	meters	.22	.67	2
Temp.	T_{e}	keV	0.1	2.0	8.
Lundquist no.	S		1×10 ⁴	7×10 ⁶	5×10 ⁸
Mode growth time	$\tau_A S^{1/2}$	s	2×10-4	9×10 ⁻³	7×10 ⁻²
Layer thickness	aS ^{-1/2}	m	2×10 ⁻³	2×10-4	8×10 ⁻⁵
zones	$N_R\!\!\times\!\!N_\theta\!\!\times\!\!N_\varphi$		3×10 ⁶	5×10 ¹⁰	3×10 ¹³
CFL timestep	$\Delta X/V_A$ (Explicit)	S	2×10 ⁻⁹	8×10 ⁻¹¹	7×10 ⁻¹²
Space- time pts			6×10 ¹²	1×10 ²⁰	6×10 ²⁴



International
Thermonuclear
Experimental
Reactor

in Cadaraches, France, operational by 2017



Slides from David Keys --Altered some for NERSC6 10¹² needed (explicit 22uniform baseline)



How to Increase Efficiency?

- In tree ed processor speed and efficiency
- Increased concurrency
- Higher-order discretizations
 - Same accuracy can be achieved with many fewer elements
- Flux-surface following gridding
 - Less resolution required along than across field lines
- Adaptive gridding
 - Zones requiring refinement are <1% of ITER volume and resolution requirements away from them are ~10² less severe
- Implicit solvers
 - Mode growth time 9 orders longer than Alfven-limited CFL







Illustrations from Computational MHD

- M3D code (Princeton)
 - multigrid replaces block Jacobi/ASM preconditioner for optimality
 - new algorithm callable across Ax=b interface
- NIMROD code (General Atomics)
 - direct elimination replaces PCG solver for robustness
 - scalable implementation of old algorithm for Ax=b



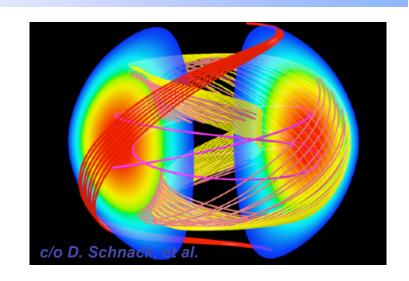




Computational MHD

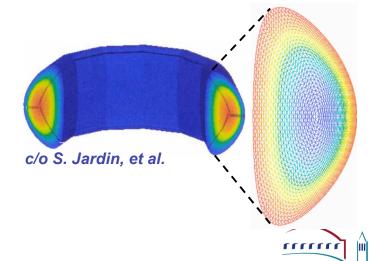
NIMROD code: Direct Elim. for robustness

- Fourier transforms in toroidal direction
- High-order finite elements in 2D poloidal crossplanes
- Sequence of complex, nonsymmetric linear systems with 10K-100K unknowns in 2D (>90% exe. Time)
- Uses SuperLU (parallel sparse direct solver benefits from efficient support of very small messages sizes)



M3D code: multigrid for optimality

- Finite differences in toroidal direction
- Unstructured mesh, hybrid FE/FD discretization with C0 elements in 2D poloidal crossplanes
- Sequence of real scalar systems (>90% exe. Time)
- algebraic multigrid (AMG) from Hypre (multigrid benefits from good support of Office of fine-grained

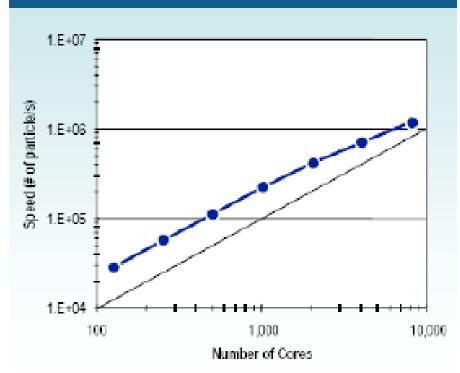






Scaling of PIC Codes

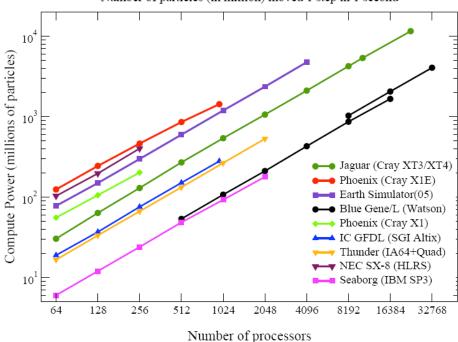
XGC Strong Scaling : 131M ions and electrons, 200K grid



FSP example (C.S. Chang)

SciDAC example (S. Ethier)

Compute Power of the Gyrokinetic Toroidal Code Number of particles (in million) moved 1 step in 1 second



S. Ethier, PPPL, Apr. 2007

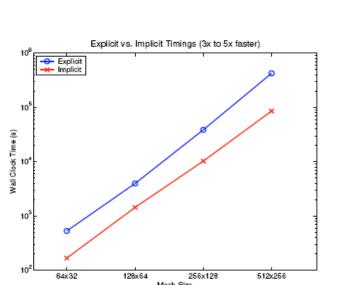


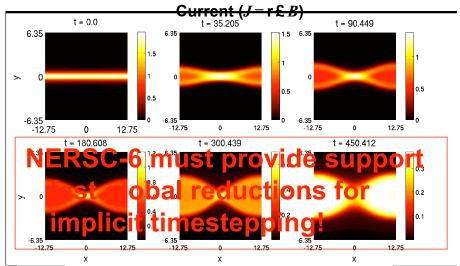


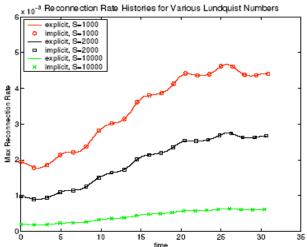


Resistive MHD: Nonlinear Implicit Model

- Magnetic reconnection: the breaking and reconnecting of oppositely directed magnetic field lines in a plasma, replacing hot plasma core with cool plasma, halting the fusion process
- Replace explicit timestepping with implicit Newton-Krylov from SUNDIALS with factor of ~5× in execution time







J. Brin et al., "Geospace Environmental Modeling (GEM) magnetic reconnection challenge," J. Geophys. Res. 106 (2001) 3/15-3/19-





Memory Bandwidth and Interconnect

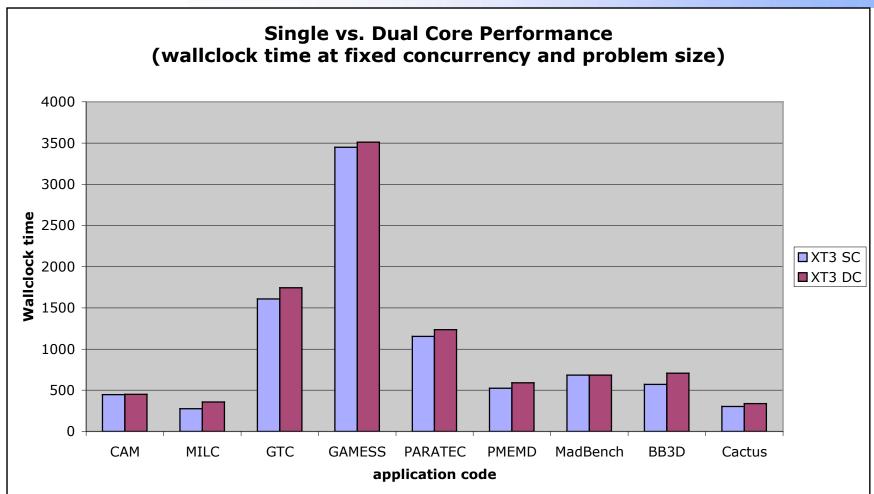






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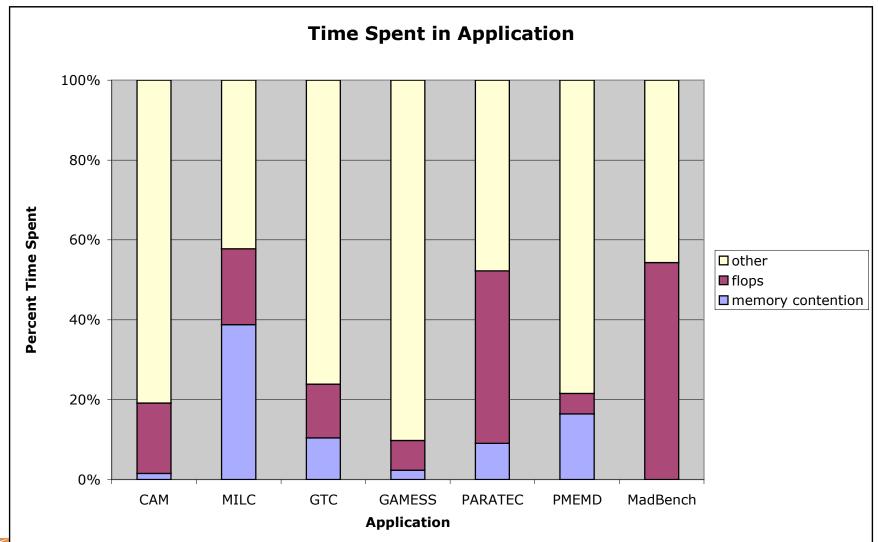
Sensitivity to Memory Bandwidth



Poor compiler performance makes applications underutilize mem bandwidth Result: relatively insensitive to halving memory bandwidth



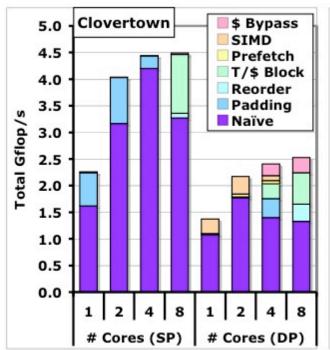
Time spent dominated by "other" (mostly latency stalls)

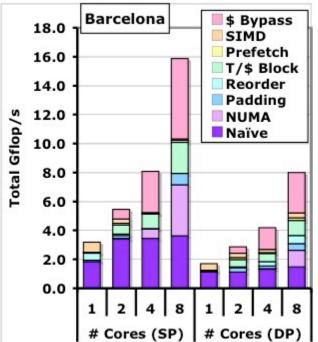


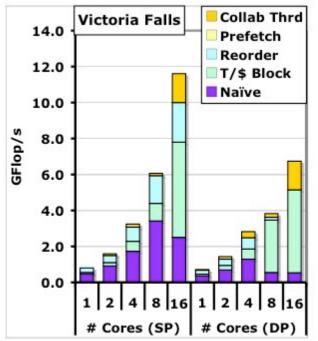




Hand-Tuned Kernels Can Reach Peak (and BW ceiling)









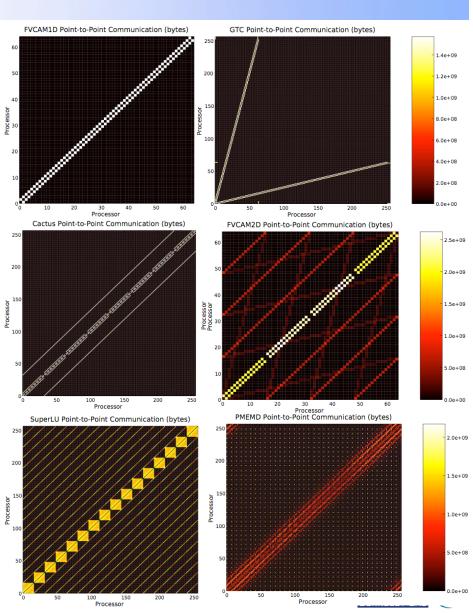




Interconnect Design Considerations for Massive Concurrency

- Application studies provide insight to requirements for Interconnects (both onchip and off-chip)
 - On-chip interconnect is 2D planar (crossbar won't scale!)
 - Sparse connectivity for dwarfs; crossbar is overkill
 - No single best topology
- A Bandwidth-oriented network for data
 - Most point-to-point message exhibit sparse topology & bandwidth bound
- Separate Latency-oriented network for collectives
 - E.g., Thinking Machines CM-5, Cray T3D, IBM BlueGene/L&P
- Ultimately, need to be aware of the on-chip interconnect topology in addition to the offchip topology
 - Adaptive topology interconnects (HFAST)
 - Intelligent task migration?







Interconnects Need For High Bisection Bandwidth

3D FFT easy-to-identify as needing high bisection

- Each processor must send messages to all PE's! (all-to-all) for 1D decomposition
- However, most implementations are currently limited by overhead of sending small messages
- 2D domain decomposition (required for high concurrency) actually requires sqrt(N) communicating partners (someto-some)

Same Deal for AMR

 AMR communication is sparse, but by no means is it bisection bandwidth limited

